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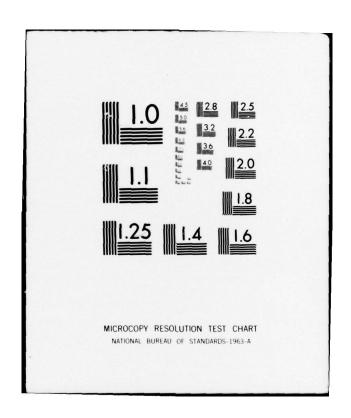






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TWO RESULTS CONCERNING MULTICOLORING

by

V. Chvatal, M. R. Garey and D. S. Johnson

STAN-CS-76-582 DECEMBER 1976

COMPUTER SCIENCE DEPARTMENT School of Humanities and Sciences STANFORD UNIVERSITY

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ABSTRACT

The m-chromatic number $\chi_m(G)$ of a graph G=(V,E) is the least integer k such that there exists a mapping $f:V+\{S\subseteq\{1,2,\ldots,k\}:|S|=m\}$ having the property that $f(u)\cap f(v)=\phi$ whenever $\{u,v\}\in E$. This is a generalization of the standard notion of chromatic number and arises in connection with mobile telephone frequency assignments. Answering a question of Lovász, our first result shows that for any $m\geq 1$ and any $\epsilon>0$, there exists a graph G for which $\chi_{m+1}(G)/\chi_m(G)>2-\epsilon$. This shows that the known bound of 2 for all m and G is essentially best possible. Our second result shows that the least integer m_0 for which $\chi_{m_0}(G)/m_0=1$ m_0 for all m and m can be asymptotically as large as m for some m vertex graphs, though it can never exceed m for some m vertex graphs, though it can never exceed m for m for m for m for some m vertex graphs, though it can never exceed m for m for m for m for some m vertex graphs, though it can never exceed m for m for m for m for m for m for some m vertex graphs, though it can never exceed m for m

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I. INTRODUCTION

The following generalization of the standard notion of graph coloring has been of recent interest [1,3,4,6,7,8]. A multicoloring of a graph G = (V,E) is a function f defined on V whose values are sets (of "colors") satisfying $f(u) \cap f(v) = \phi$ whenever $\{u,v\}$ ϵ E. For positive integers k, m, a (k,m)-coloring of G = (V,E) is a multicoloring f of G such that |f(v)| = m for each v ϵ V and $|\bigcup f(v)| = k$. The m-chromatic number $\chi_m(G)$ is the least integer k such that there exists a (k,m)-coloring of G. (This last definition differs from that of [6,7] by a factor of m.) Notice that for m = 1 these definitions correspond to the usual graph coloring notions. The purpose of this note is to resolve two questions about multicoloring conveyed to us by P. Erdös [2].

The first question deals with the relationship between $\chi_m(G)$ and $\chi_{m+1}(G)$. It is not difficult to see that

 $\chi_{m+1}(G) \leq \chi_m(G) + \chi_1(G) \leq 2 \cdot \chi_m(G)$



with equality possible in the right-hand inequality only for m = 1. Lovász asked [2] whether, for each value of m, there exist graphs G such that $\chi_{m+1}(G) > (2-\epsilon)\chi_m(G)$. We shall answer this question in the affirmative.

The graphs we shall use are defined as follows: for positive integers $n\geq 2m$, the graph G^n_m has vertex set consisting of all m-element subsets of $\{1,2,\ldots,n\}$ and has an edge between two such vertices exactly when their intersection is empty. It is easy to see that $\chi_m \left(G^n_m\right) \leq n$ merely by considering the multicoloring provided by the definition of G^n_m and, in fact, it is proved in [7,8] that $\chi_m \left(G^n_m\right) = n$. Thus, to answer the question of Lovász, it suffices to prove the following theorem:

Theorem 1. For each $m \ge 2$, there exists a constant c such that for all sufficiently large n

$$\chi_{m+1}(G_m^n) \ge 2n - c.$$

In order to prove Theorem 1, we require the following lemma, which is an immediate consequence of a special case of Theorem 3 in [5].

<u>Lemma 1</u>. For fixed $m \ge 2$ and n sufficiently large, there exists a constant a_0 such that the number of m-element subsets of $\{1,2,\ldots,n\}$ which can be chosen so that no two are disjoint but there is no element common to all is at most a_0^{m-2} .

<u>Proof of Theorem 1</u>. Fix m. We merely need to show that, for all sufficiently large n,

$$\chi_{m+1}(G_m^n) - \chi_{m+1}(G_m^{n-1}) \ge 2$$

and the result will follow by induction. So suppose we have a (k,m+1)-coloring of G_m^n such that $k=\chi_{m+1}\left(G_m^n\right)$, where n is any integer sufficiently large that the conclusion of Lemma 1 holds and such that $\binom{n-1}{m-1} > ma_0 n^{m-2}$, where a_0 is the constant of Lemma 1. We first claim that there must be at least n+1 colors which each appear on more than $a_0 n^{m-2}$ vertices.

Suppose there are n or fewer colors which each appear on more than $a_0 n^{m-2}$ vertices. By Lemma 1, each such color can appear on at most $\binom{n-1}{m-1} > a_0 n^{m-2}$ vertices since they must all share a common element. Thus, since each of the $\binom{n}{m}$ vertices receives exactly m+1 colors, we must have

$$(m+1) \quad {n \choose m} \le n \quad {n-1 \choose m-1} + (k-n)a_0 n^{m-2}$$

$$\le m \quad {n \choose m} + (k-n)a_0 n^{m-2}$$

or, rewriting,

$$\binom{n}{m} \leq (k-n)a_0 n^{m-2}$$

Since

$$k = \chi_{m+1} (G_m^n) \le \chi_m (G_m^n) + \chi (G_m^n)$$

$$\le 2\chi_m (G_m^n) = 2n$$

it follows that we must have

$$\binom{n}{m} \leq a_0 n^{m-1}$$

However this is a contradiction, since n was chosen sufficiently large that $\binom{n}{m} = \frac{n}{m} \binom{n-1}{m-1} > \frac{n}{m} \max_{0} n^{m-2} = a_{0} n^{m-1}$, and the claim follows.

Thus there are at least n+l colors which each appear on more than $a_0 n^{m-2}$ vertices. The set of vertices on which any color i appears must form a collection of pairwise-intersecting m-element subsets of $\{1,2,\ldots,n\}$, by definition of G_m^n . Thus, by Lemma 1, whenever color i appears on more than $a_0 n^{m-2}$ vertices, all those vertices must contain some common element e_i . Since there are more than n such colors, we must have $e_i = e_j$ for some i and j. If we delete from G_m^n all the vertices containing $e_i = e_j$, we obtain a copy of G_m^{n-1} and a (k-2,m+1)-coloring of it, since colors i and j have disappeared. Therefore

$$\chi_{m+1} (G_m^{n-1}) \le k-2 = \chi_{m+1} (G_m^n) - 2$$

and the theorem is proved.

The second question involves what we call the ultimate multichromatic number $\chi^*(G)$ defined by

$$\chi^*(G) = \inf_{m} \chi_m(G)/m.$$

It is proved in [1,7] that the value of $\chi^*(G)$ is always achieved for some finite m. One easy way to see this is to formulate the problem of determining $\chi^*(G)$ as a linear programming problem (as done in [4]): Let v_1, v_2, \ldots, v_n be an ordering of the vertices of G and let S_1, S_2, \ldots, S_ℓ be an ordering of the independent sets of G. Define x_{ij} to be 1 whenever $v_i \in S_j$ and 0 otherwise. Then the value of $\chi^*(G)$ is given by

$$\chi^*(G) = \min \sum_{j=1}^{\ell} r_j$$

subject to: $r_j \ge 0$, $1 \le j \le l$;

$$\sum_{j=1}^{\ell} x_{ij}r_j = 1, \quad 1 \leq i \leq n.$$

One can show easily, using Hadamard's Theorem, that no basis matrix for this problem can have determinant exceeding $n^{n/2}$ and this is an upper bound on the value of m required.

This upper bound however seems ridiculously large. Erdös asked [2] (as did the authors, independently) whether $\chi^*(G)$ could always be achieved for an m not exceeding the number of vertices of G. We answer this in the negative, constructing graphs for which extremely large values of m are necessary to achieve $\chi^*(G)$.

Let C_p denote the graph which is a cycle on p vertices. The join G_1+G_2 of two graphs G_1 and G_2 , having disjoint vertex sets, consists of all edges and vertices in the two given graphs along with all edges joining a vertex from G_1 to a vertex from G_2 . We use the following two lemmas in our construction:

Lemma 2. [4,7] For all integers $p \ge 1$,

$$\chi^*(C_{2p+1}) = 2 + (1/p).$$

Lemma 3. [7] For all graphs G_1 and G_2 ,

$$\chi^*(G_1 + G_2) \; = \; \chi^*(G_1) \; + \; \chi^*(G_2).$$

Let p_i denote the ith prime and define the graph G(i) to be $C_{2p_1+1} + C_{2p_2+1} + \cdots + C_{2p_i+1}$. The number of vertices n of G(i) is given by

$$n = i + 2 \sum_{j=1}^{i} p_{j}$$
.

Applying Lemmas 2 and 3, we obtain

$$\chi^*(G(1)) = 2i + \sum_{j=1}^{1} (1/p_j)$$
.

Since $\chi_m(G(i))$ must always be an integer, it follows that the least value of m for which $\chi^*(G(i)) = \chi_m(G(i))/m$ can be no less that $\prod_{j=1}^{n} p_j$ (and in fact that value of m will work). Using the Prime Number Theorem and expressing this lower bound in

terms of n, we obtain the asymptotic lower bound of

 $e^{\sqrt{(n\log n)/2}}$.

Thus, though this is still quite far from the upper bound of

 $n^{n/2} = e^{(n\log n)/2}$

we see that extremely large values of m $\underline{\text{can}}$ be required in order to achieve $\chi^*(G)$.

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